

## LOWER BOUNDS FOR CORNER-FREE SETS

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Abstract. A *corner* is a set of three points in  $\mathbf{Z}^2$  of the form  $(x, y), (x + d, y), (x, y + d)$  with  $d \neq 0$ . We show that for infinitely many  $N$  there is a set  $A \subset [N]^2$  of size  $2^{-(c+o(1))\sqrt{\log_2 N}} N^2$  not containing any corner, where  $c = 2\sqrt{2\log_2 \frac{4}{3}} \approx 1.822\dots$

Let  $q, d$  be large positive integers. For each  $x \in [q^d - 1]$ , we may write  $\pi(x) = (x_0, \dots, x_{d-1}) \in \mathbf{Z}^d$  for the vector of digits of its base  $q$  expansion, thus  $x = \sum_{i=0}^{d-1} x_i q^i$ , with  $0 \leq x_i < q$  for all  $i$ .

For each positive integer  $r$ , consider the set  $A_r$  of all pairs  $(x, y) \in [q^d - 1]^2$  for which  $\|\pi(x) - \pi(y)\|_2^2 = r$  and  $\frac{q}{2} \leq x_i + y_i < \frac{3q}{2}$  for all  $i$ .

We claim that  $A_r$  is free of corners. Suppose that  $(x, y), (x + d, y), (x, y + d) \in A_r$ . Then

$$\|\pi(x) - \pi(y)\|_2^2 = \|\pi(x + d) - \pi(y)\|_2^2 = \|\pi(x) - \pi(y + d)\|_2^2 = r. \quad (1)$$

We claim that

$$\pi(x + d) + \pi(y) = \pi(x) + \pi(y + d). \quad (2)$$

To this end, we show that  $(x + d)_i + y_i = x_i + (y + d)_i$  for  $i = 0, 1, \dots$  by induction on  $i$ . A single argument works for both the base case  $i = 0$  and the inductive step. Suppose that, for some  $j \geq 0$ , we have the statement for  $i < j$ . Write  $x_{\geq j} := \sum_{i \geq j} x_i q^i$ , and define  $(x + d)_{\geq j}, y_{\geq j}, (y + d)_{\geq j}$  similarly. By the inductive hypothesis and the fact that  $x + (y + d) = (x + d) + y$ , we see that  $x_{\geq j} + (y + d)_{\geq j} = (x + d)_{\geq j} + y_{\geq j}$ . Therefore  $x_j + (y + d)_j = (x + d)_j + y_j \pmod{q}$ . However by assumption we have  $\frac{q}{2} \leq x_j + (y + d)_j, (x + d)_j + y_j < \frac{3q}{2}$ , and so  $x_j + (y + d)_j = (x + d)_j + y_j$ . The induction goes through.

With (2) established, let us return to (1). We now see that this statement implies that  $\|a\|_2^2 = \|a + b\|_2^2 = \|a - b\|_2^2 = r$ , where  $a := \pi(x) - \pi(y)$  and  $b := \pi(x + d) - \pi(x) = \pi(y + d) - \pi(y)$ . By the parallelogram law  $2\|a\|_2^2 + 2\|b\|_2^2 = \|a - b\|_2^2 + \|a + b\|_2^2$ , this immediately implies that  $b = 0$ . Since  $\pi$  is injective, it follows that  $d = 0$  and so indeed  $A_r$  is corner-free.

The set of all pairs  $(x, y)$  with  $\frac{q}{2} \leq x_i + y_i < \frac{3q}{2}$  for all  $i$  has size  $(\frac{3}{4}q^2 + O(q))^d$ . Therefore by the pigeonhole principle there is some  $r$  such that  $\#A_r \geq (dq^2)^{-1}(\frac{3}{4}q^2 + O(q))^d$ .

Now for a given  $d$  set  $q := \lfloor (2/\sqrt{3})^d \rfloor$  and  $N := q^d$ . Then  $A_r \subset [N] \times [N]$ ,  $A_r$  is free of corners, and

$$\#A_r \geq N^2(dq^2)^{-1} \left( \frac{3}{4} + O\left(\frac{1}{q}\right) \right)^d.$$

Writing  $o(1)$  for a quantity tending to 0 as  $N \rightarrow \infty$ , we note that  $q = (\frac{2}{\sqrt{3}} + o(1))^d$  and that  $d = (1 + o(1)) \sqrt{\frac{\log_2 N}{\log_2(2/\sqrt{3})}}$ . A short calculation then confirms that

$$\#A_r \geq N^2 2^{-(c+o(1))\sqrt{\log_2 N}},$$

where  $c = 2\sqrt{2 \log_2 \frac{4}{3}} \approx 1.822 \dots$

*Remark.* The construction came about by a careful study of the recent preprint of Linial and Shraibman [1], where they used ideas from communication complexity to obtain a bound with  $c = 2\sqrt{\log_2 e} \approx 2.402 \dots$ , improving on the previously best known bound with  $c = 2\sqrt{2} \approx 2.828 \dots$  which comes from Behrend's construction. By bypassing the language of communication complexity one may simplify the construction, in particular avoiding the use of entropy methods. This yields a superior bound.

## References

- [1] N. Linial and A. Shraibman, *Larger corner-free sets from better NOF exactly- $N$  protocols*, preprint, arxiv:2102.00421.

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